

#### Efficiently Querying Protein Sequences with the Proteinus Index

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#### Contents

- Motivation
- Background
  - Reduced Alphabet
  - The nsP-index
- The Proteinus
- Experiments
- Conclusions



### Motivation

- Finding similarities in protein sequences stored in biological databases has become a core problem in bioinformatics
- Applications
  - Protein evolution studies
  - Predicting biological structure
  - Functional characterization of novel protein sequences
  - ...
  - And much more...



# Protein sequence databases

- Protein sequence databases (PSD) store data related to the primary structure of proteins and their amino acids chains
  - These chains are stored in the databases as strings
    - Representing the 20 molecules present in proteins

 $\Sigma_{20}^{p} = \{A, C, D, E, F, G, H, I, K, L, M, N, P, Q, R, S, T, V, W, Y\}$ 

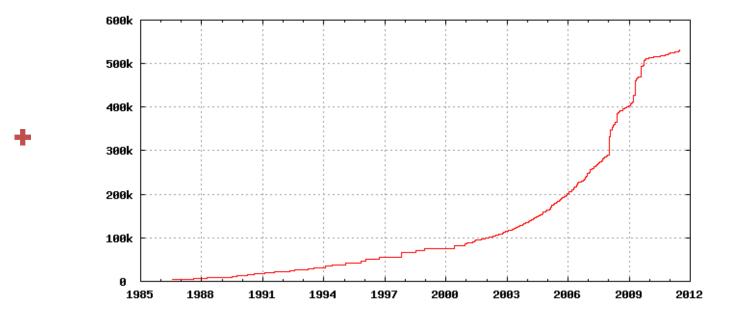


### Similarities in PSD

- An important type of query used to search for similarities in PSD is the range query
- The Edit Distance (ED) is the most used distance function to compare sequences
  - ED is defined as the minimum number of edit operations
  - (insertion, deletion and update)
- Due to protein sequences characteristics, the ED function must consider weights related to amino acid exchanges
  - these weights are obtained by *replacing matrices*, such as the PAM and the BLOSUM



With the technological advances in molecular biology, the amount of protein sequences has increased exponentially



Number of entries in UniProtKB/Swiss-Prot



#### Similarities in PSD

• In the literature, tools widely used to search for similarities in PSD are:

– BLASTP, PSI-BLAST and PHI-BLAST +

Another research area uses indices to improves the performance of similarity search

– MRS, ComRI, nsP-Index

However, these approaches are aimed at indexing nucleotide sequences databases, and are not directly applicable on PSD



### Similarities in PSD

- PSD are different from nucleotide sequences databases
  - PSDs store a greater number of protein sequences
  - Protein sequences (typically of hundreds to thousands residues) are much smaller than nucleotide sequences (typically of millions of bases)
  - The amino acid alphabet is bigger than the nucleotide alphabet (A, C, G, T)

Recall that computing similarity among protein sequences is necessary to handle *replacing matrices* (e.g., PAM, BLOSUM, ...)



#### Proposal

- We propose Proteinus index (acronym for Protein sEquence sImilarity Search)
  - a new protein sequence persistent index
  - It uses a reduced amino acid alphabet
    proposed by Li *et al.* [1]
  - It allows the persistent storage of the index on disk
    - Proteinus extends the nsP-index

[1] Li, T., Fan, K., Wang, J., Wang, W.: Reduction of protein sequence complexity by residue grouping. Protein Eng. 16(5), 323-330 (2003)



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# Reduced Amino Acid Alphabet

- Some amino acids (e.g., ILE and LEU) are chemically similar, and can be replaced by one another without changing the global structure (the fold) or the function of a protein
- The amino acid alphabet can be reduced to a new alphabet wherein similar amino acids are clustered together



- The first reduced amino acid alphabet was introduced by Dill [2]:
  - Hydrophobic-polar (HP) model for study of the folding of globular proteins
- Since then, more than 50 reduced alphabets of different sizes have been proposed

[2] Dill K.: "Theory for the Folding and Stability of Globular Proteins" Biochemistry, 24:1501-1509, 1985

### Reduced Amino Acid Alphabet

- Li et al [3] proposed and compared several reductions on the complete alphabet
  - by grouping amino acids based on their physical an chemical characteristics
  - by seeking to maximize a similarity score derived from the BLOSUM62 matrix
- The authors concluded that a reduced alphabet, which is composed by 10 groups, preserves the maximum information on the original protein sequence

Using a reduced alphabet in protein sequence comparisons avoids using a replacing matrix



$$\Sigma_{20}^p = \{A, C, D, E, F, G, H, I, K, L, M, N, P, Q, R, S, T, V, W, Y\}$$

 $\Sigma_{10}^{p} = \{C', F', M', I', G', P', A', N', Q', R'\}, such that$ 

$$\begin{split} C' &= \{C\}, \ F' = \{F, Y, W\}, \ M' = \{M, L\}, \ I' = \{I, V\}, \ G' = \{G\}, \ P' = \{P\}, \\ A' &= \{A, T, S\}, \ N' = \{N, H\}, \ Q' = \{Q, E, D\}, \ R' = \{R, K\} \end{split}$$



#### Contents

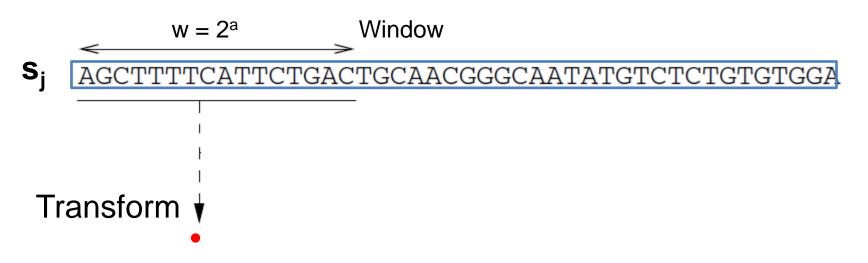
- Motivation
- Background
  - Reduced Alphabet
  - The nsP-index
- The Proteinus
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- The nsP-Index extends the MRS-index into a persistent index aimed at nucleotide sequences
  - Let S<sub>n</sub> = {s1, s2, ..., s<sub>x</sub>} be a database composed of x nucleotide sequences
  - and w = 2<sup>a</sup> be the length of the shortest query sequence

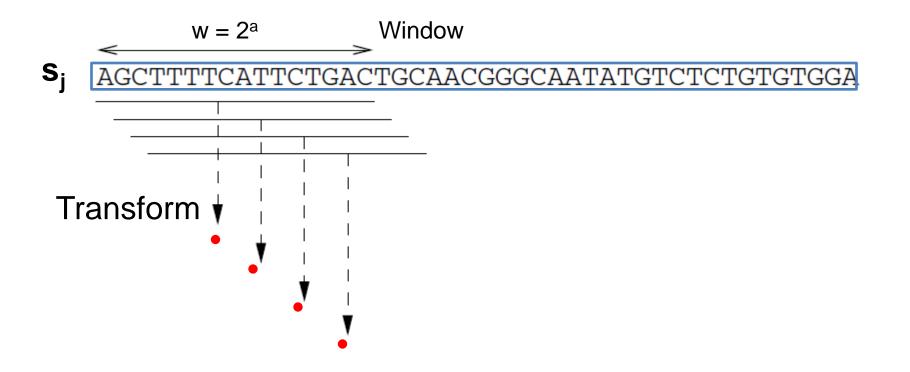


- For all sequence s<sub>i</sub> in the database
  - A window of length  $w = 2^i$  ( $a \le i \le b$ ) slides over the sequence s<sub>i</sub>
- For each placement of the window, a subsequence of s<sub>j</sub> is mapped into a point using a wavelet-based method



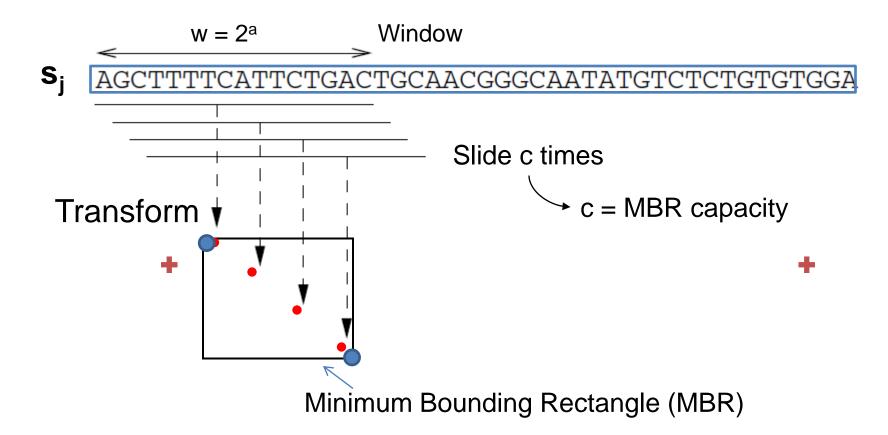


• The sliding window starts from the leftmost point to the rightmost point of the sequence



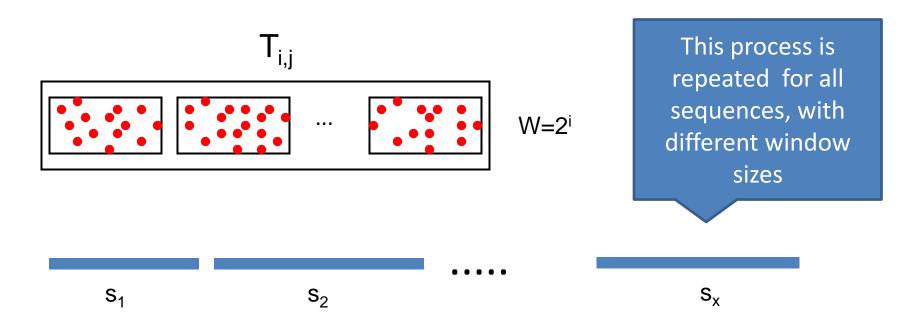


• These points are comprised within a minimum bounding rectangle (MBR)





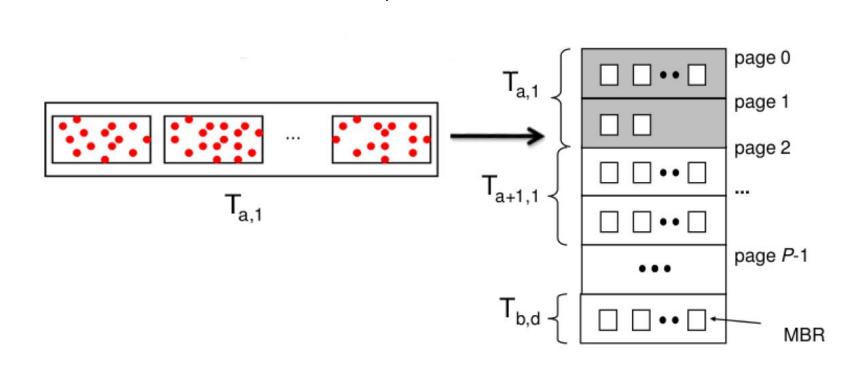
- A list of MBRs represents an element T<sub>i,i</sub>
- Each element T<sub>i,j</sub> indexes the j<sup>th</sup> sequence with window sizes 2<sup>i</sup>



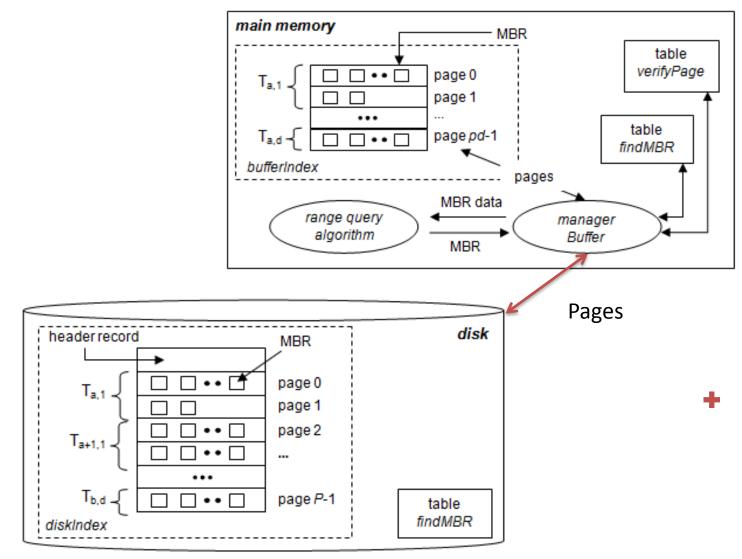


These T<sub>i,j</sub> elements are stored on disk, in different pages

S<sub>i</sub>



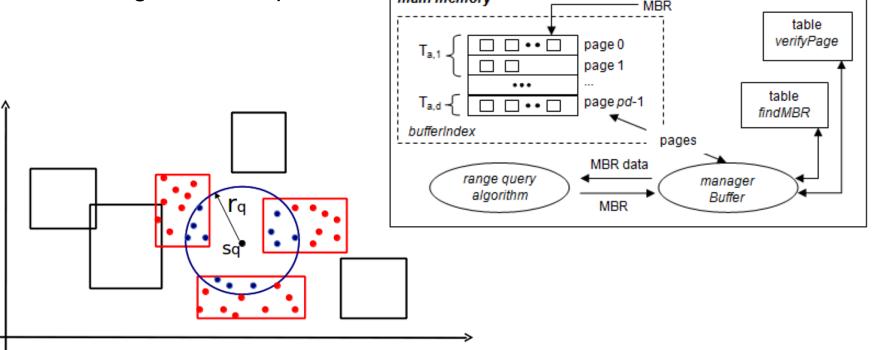
## The nsP-index (physical view)



### The nsP-index (range query)

 Range query processing calculates the distance between the (transformed point) query sequence and the MBRs

using the buffer-pool



main memory

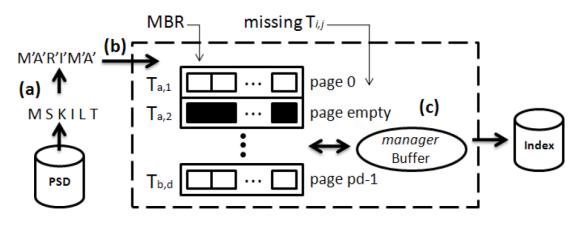


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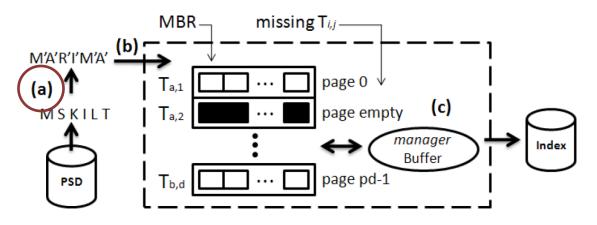
- The Proteinus Index extends the nsP-Index to process protein sequences
  - Inherits characteristics of the nsP-Index
  - Uses a reduced amino acid alphabet



**Overview of the Proteinus approach** 



- The Proteinus Index extends the nsP-Index to carry on protein sequences
  - Proteinus inherits characteristics of the nsP-Index
  - It uses a reduced amino acid alphabet

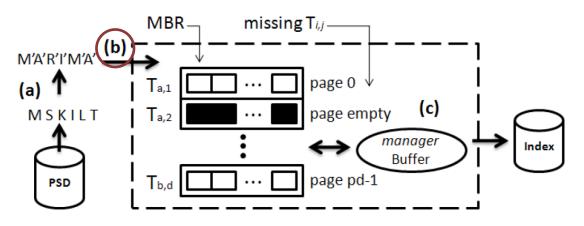


Overview of the Proteinus approach

A protein sequence is extracted from the PSD



- The Proteinus Index extends the nsP-Index to carry on protein sequences
  - Proteinus inherits characteristics of the nsP-Index
  - It uses a reduced amino acid alphabet

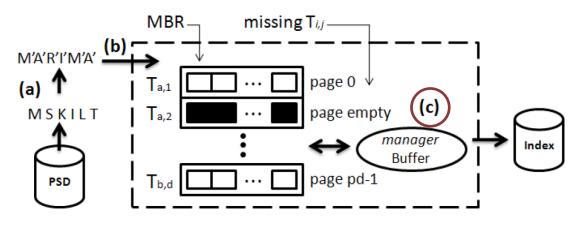


**Overview of the Proteinus approach** 

The extracted protein sequence is converted following a reduced amino acid alphabet



- The Proteinus Index extends the nsP-Index to carry on protein sequences
  - Proteinus inherits characteristics of the nsP-Index
  - It uses a reduced amino acid alphabet



**Overview of the Proteinus approach** 

The converted protein sequence is then indexed by Proteinus



- Proteinus focuses on five tasks:
  - 1. It extends the index data structure to handle the protein sequence alphabet
  - 2. It defines a mapping function to transform a subsequence s<sub>j</sub> into a multidimensional point
  - 3. It defines an approximated distance function among the transformed points
  - 4. It adapts the index structure to handle missing  $T_{i,i}$
  - 5. It provides a range query routine to support the new characteristics of the index data structure





 To accomplish Task 1, we designed a dedicated Proteinus data structure to support the reduced alphabet:

 $\Sigma_{10}^p = \{C', F', M', I', G', P', A', N', Q', R'\}$ 

- Maintaining the same memory requirements of the nsP-Index:
  - Through the use of a shorter data type to represent a protein sequence



- Proteinus's data structure increases twice the storage requirements
  - as the reduced amino acid alphabet has 10 symbols instead of 5 symbols to handle nucleotide sequences.
- Due to the small sizes of the protein sequences
  - Proteinus-index decreases by half the storage requirements through the use of a shorter data type.

Therefore, we used a 16 bit variable in the Proteinus Index to represent each multidimensional point, while the nsP-index uses a 32 bit variable



- Regarding tasks 2 and 3,
  - The mapping and approximated distance functions should be easily computable
  - To do so, we used:
    - The wavelet-method inherited from the nsP-Index
    - Together with the reduced alphabet

Defines a mapping function to transform a subsequence sj into a multidimensional point
 3. Defines an approximated distance function among the transformed points



- Task 4 is required since the protein sequence sizes vary greatly
  - Thus, not all sequences can be indexed in all resolutions
  - To this end, Proteinus data structure allow missing T<sub>i,i</sub>
  - When a sequence  $s_i$  is shorter than a window resolution, an empty  $T_{i,i}$  is generated
    - Empty Ti,j do not persist on disk



- To accomplish Task 5, we extended the nsP-index range query algorithm to support missing T<sub>i,j</sub>
  - When a query sequence  $\boldsymbol{s}_q$  is larger than a sequence  $\boldsymbol{s}_i$  in PSD
    - Our search algorithm bypasses sj in the similarity evaluation task

5. It provides a range query routine to support the new characteristics of the index data structure



- Finally, the use of approximations, such as MBRs, introduces loss of accuracy in the exact representation
  - Proteinus returns candidates (in MBRs)
  - Thus, Proteinus should provide both a filter phase and a refinement phase
  - To this end, we also developed a simple post-processing routine to eliminate false answers
    - Smith-Waterman algorithm





#### Contents

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# Setup of experiments

- The experiments were conducted on a computer with:
  - Intel Core i7 2.67 GHz processor,
  - 12 GB of main memory,
  - 2 SATA 1 TB hard disks and
  - Linux Ubuntu 9.04.
  - We employed the BLASTP version 1:2.2.21.20090809-1.
- Proteinus was implemented in C++
- We used four resolutions with window sizes of 32, 64, 128 and 256 characters, and set the MBR capacity to 80 strings
- The bufferIndex size was set to 200 MB and the disk page size c was set to 672 bytes.
  - These values were identified experimentally



# Setup of experiments

- In our tests, we used real-world protein sequences obtained from the Uniprot database
  - We created several PSDs containing increasing volumes of protein sequences: 4 GB, 6 GB, 8 GB (original dataset) and 10 GB (synthetic dataset)
- We perform range queries by
  - submitting four different sizes of queries to each PSD
    - 128, 256, 512 and 1,024 amino acids
  - With an error rate E = 0.005 to calculate query radius





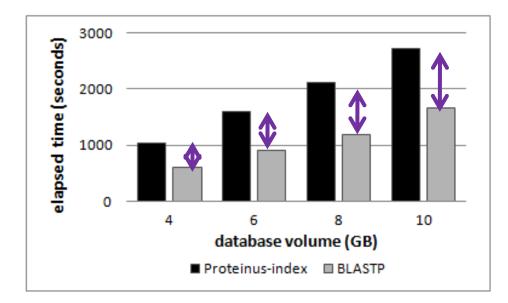
# Setup of experiments

- We compared the performance of Proteinus with the BLASTP tool
  - Taking into account the elapsed time (in seconds)
- Although Proteinus is not a complete alignment strategy as the BLASTP.....
- to have a measure of comparison



#### Experiment 1 — Index Construction

- Time spent by Proteinus to build the indices
  - Compared to the time spent by BLASTP to create its data structure (command *formatdb*)

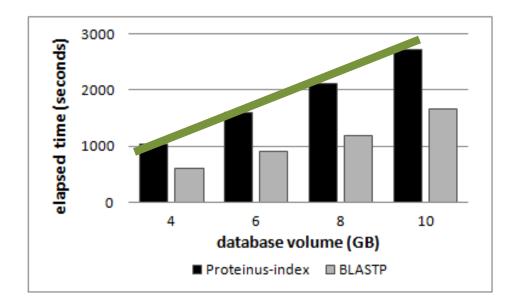


The time spent by Proteinus was (on average 43%) higher than that spent by BLASTP



#### Experiment 1 — Index Construction

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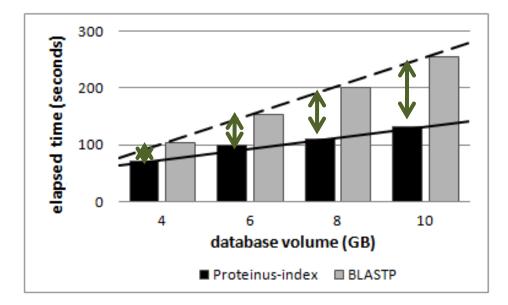


#### However, Proteinus proved to be scalable



#### Experiment 2 – Range Query

 Proteinus provided an impressive performance gain that ranged from 45% up to 93% over the BLASTP

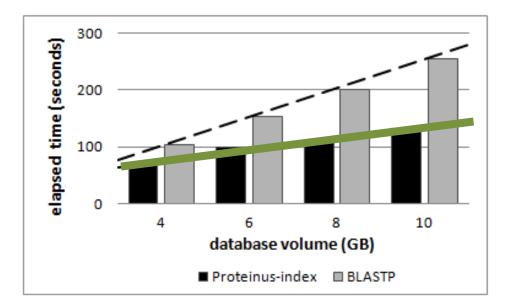


The larger the dataset, the better the performance gain



#### Experiment 2 – Range Query

 Proteinus provided an impressive performance gain that ranged from 45% up to 93% over the BLASTP

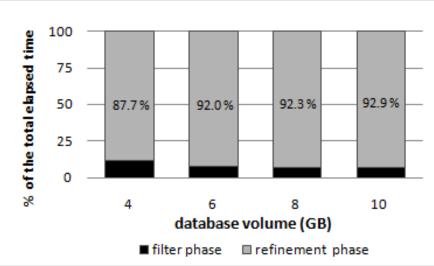


Also, Proteinus proved to be scalable in query processing



### Experiment 3 – Range Query

 Percentage of the total elapsed time spent by Proteinus to process the filter and the refinement phases separately

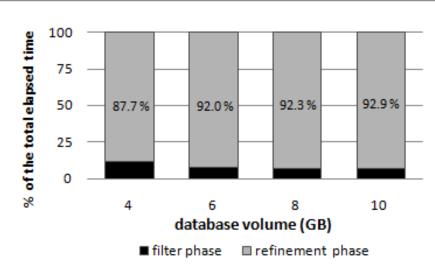






### Experiment 3 – Range Query

 Percentage of the total elapsed time spent by Proteinus to process the filter and the refinement phases separately





The larger the dataset, the cost of filter phase is maintained



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- Background
  - Reduced Alphabet
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# Conclusions

- We presented Proteinus, a new index based on approximations to process range queries over PSD
- The results showed that Proteinus was efficient in query processing, providing an performance gain that range from 45% up to 93% over BLASTP tool

Despite the time need for index creation

- Proteinus proved to be scalable both
  - Index creation
  - Query processing



## Conclusions

+ Proteinus can use any reduced amino acid alphabet

- Proteinus can be used as an efficient filter phase to select a reduced set of protein sequences
  - which can be further analyzed by any post-processing tool





## Questions? Suggestions?

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# Similarities in PSD

- An important type of query used to search for similarities in PSD is the range query
  - Consider that S = {s1, s2, ..., sd} is a PSD
  - Sq represents a protein sequence query
  - And d(si, sq) represents the distance between a protein sequence si stored in S and sq
  - Given a query radios rq, the range query retrieve every protein sequence that satisfies the condition d(si,sq) <= rq.</li>

